Non-Refined Abstractions in Counterexample Guided Abstraction Refinement for Multi-Agent Path Finding (Extended Abstract)

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Abstract
Counterexample guided abstraction refinement (CEGAR) represents a powerful symbolic technique for various tasks such as model checking and reachability analysis. Recently, CEGAR combined with Boolean satisfiability (SAT) has been applied for multi-agent path finding (MAPF), a problem where the task is to navigate agents from their start positions to given individual goal positions so that agents do not collide with each other. The recent CEGAR approach used the initial abstraction of the MAPF problem where collisions between agents were omitted and were eliminated in subsequent abstraction refinements. We propose in this work a novel CEGAR-style solver for MAPF based on SAT in which some abstractions are deliberately left non-refined. This adds the necessity to post-process the answers obtained from the underlying SAT solver as these answers slightly differ from the correct MAPF solutions. Non-refining however yields order-of-magnitude smaller SAT encodings than those of the previous approach and speeds up the overall solving process.

Introduction
Multi-agent path finding (MAPF) (Silver 2005) is a task of finding non-conflicting paths for $k \in \mathbb{N}$ agents $A = \{a_1, a_2, \ldots, a_k\}$ that move in an undirected graph $G = (V, E)$ across its edges such that each agent reaches its goal vertex from the given start vertex via its path. Starting configuration of agents is defined by a simple assignment $s : A \rightarrow V$ and the goal configuration is defined by a simple assignment $g : A \rightarrow V$. A conflict between agents is usually defined as simultaneous occupancy of the same vertex by two or more agents or as a traversal of an edge by agents in opposite directions.

We address MAPF from the perspective of compilation techniques that represent a major alternative to search-based solvers (Sharon et al. 2015) for MAPF. Compilation-based solvers reduce the input MAPF instance to an instance in a different well established formalism for which an efficient solver exists. Such formalisms are for example constraint programming (CSP), Boolean satisfiability (SAT), or mixed integer linear programming (MILP).

The basic compilation scheme (see Figure 1) for sum-of-costs optimal MAPF solving has been introduced by the MDD-SAT (Surynek et al. 2016) solver that uses so called complete models to compile MAPF instances into SAT. The target Boolean formula of the complete model is satisfiable if and only if the input MAPF has a solution of a specified sum-of-costs. The complete model as introduced in MDD-SAT consists of three group of constraints:

- **Agent propagation** constraints - these constraints ensure that if an agent appears in vertex $v$ at time step $t$ then the agent appears in some neighbor of $v$ (including $v$) at time step $t + 1$. The side effect of these constraints is that the agent never disappears. Cost calculation and bounding is done together with agent propagation.
- **Path consistency** constraints - these constrains ensure that agents move along proper paths, that is, agents do not clone themselves and do not appear spontaneously.
- **Conflict elimination** constraints - to ensure that agents do not conflict with each other according to the MAPF rules (vertex and edge conflict).

A significant improvement over complete models in problem compilation for MAPF is the introduction of laziness via incomplete models where the conflict elimination constraints are omitted for which equivalent solvability no longer holds, but only the implication: if the MAPF instance is solvable then the instance in the target formalism is solvable too (Surynek 2019).

The discrepancy between the original formulation of MAPF and its compiled variant in the target formalism is eliminated by abstraction refinements similarly as it is done in the counterexample guided abstraction refinement (CEGAR) (Clarke et al. 2000) approach for model checking (see Figure 2).

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CEGAR for MAPF

The general CEGAR approach for compilation-based problem solving starts with a so called initial abstraction of the problem instance being solved in some target formalism such as SAT or CSP. The initial abstraction do not model the input instance in the full details. However still the initial abstraction is passed to the solver for the target formalism despite the solver is not provided all the details needed to solve it. Then the solver will come with some answer and since it could not take some details into account during the solving phase, the answer must be checked, usually against full details of how the problem instance is defined.

Two cases need to be distinguished at this stage. If the provided answer matches the instance definition then it is returned and the solving process finishes. Otherwise the CEGAR solving process generates counterexample that is determined by the mismatch between the provided answer and the requirements the expected answer should satisfy. This mismatch is usually represented by the violation of some constraints that were not expressed in the abstraction. Then the solving process continues with a so called abstraction refinement in which the abstraction is augmented to eliminate the counterexample and the solving process continues with the next iteration of (now refined) abstraction solving (see Figure 2).

Non-Refined Abstractions in MAPF

We are going further in the CEGAR architecture of the MAPF solver. In addition to conflict elimination constraints we also omit path consistency constraints in the initial abstraction. Moreover we never make any refinement with respect to the omitted path consistency constraints - the corresponding abstraction remains non-refined.

Omitting the path consistency constraint however leads to solving a different though equisatisfiable problem. Instead of attempting to connect the initial positions of agents with their goals positions via paths, we are now attempting to make the connection via directed acyclic graphs (DAGs). This requires addition of a new polynomial-time post-processing step in the solving process that extracts valid paths from DAGs (see Figures 3 and 4).

Conclusion

Experiments show promising results as non-refining yields order of magnitude smaller formulae and faster solving runtimes than previous CEGAR-inspired SAT-based MAPF solver. We believe that non-refined abstractions in CEGAR-inspired MAPF solving opens new ways in compilation-based approaches for MAPF. Non-refining path consistency constrains is just one example and we believe that other abstractions can be discovered.

References


